# Query optimization in data stream processing

Optimizarea interogărilor în procesarea fluxurilor de date

Abstract

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## The thesis contains the following chapters<sup>1</sup>:

List of figures

List of tables

### 1 Introduction

- 1.1 Data stream processing in pervasive environments
- 1.2 Research directions
- 1.3 Original contributions
- 1.4 Thesis structure

## 2 Data stream processing. State of the art

- 2.1 Continuous processing paradigm. Overview
- 2.2 STREAM, Aurora, Medusa and Borealis
- 2.3 Conclusions

## 3 Optimizing resource usage in data stream processing

3.1 Introduction

3.2 Addressing resource usage in data stream query processing: the sizing window effect

3.3 The kSiEved Window Training Set technique

3.4 Conclusions

## 4 Resource-aware architectures for data stream processing

- 4.1 Introduction
- 4.2 An architecture for the sizing window effect in data stream processing
- 4.3 An architecture for load shedding operations in data stream processing

4.4 A solution for evaluating performance in a monitoring application with StreamInsight: StreamEval

4.5 Recommendations for data stream processing in specific application domains

4.6 Conclusions

## 5 Heterogeneous data management in a pervasive environment

5.1 Introduction

<sup>&</sup>lt;sup>1</sup>We don't expand the sections in this abstract.

- 5.2 Pervasive computing and pervasive applications. Context
- 5.3 Scenario and testbed
- 5.4 Using a system for pervasive environments in the testbed
- 5.5 Demo
- 5.6 Conclusions

6 Evaluating agility in pervasive data-oriented application development

- 6.1 Introduction
- 6.2 Data-oriented pervasive application development: used systems
- 6.3 The AgilBench benchmark
- 6.4 Evaluated systems
- 6.5 Experimental study
- 6.6 Experimental results analysis
- 6.7 AgilBench innovation
- 6.8 Conclusions
- 7 Conclusion
  - 7.1 Obtained results and research directions
  - 7.2 Final words

#### Bibliography

Keywords: data streams, continuous queries, data stream management systems, query optimization, reducing resource usage, performance optimization, pervasive applications, pervasive computing, heterogeneous data management

## Publications related to this thesis

The results of our research and the original contributions presented in the thesis were published in journals or proceedings of the international conferences we attended (one of these papers is accepted for publication):

- Sabina Surdu and Vasile-Marian Scuturici, Addressing resource usage in stream processing systems: sizing window effect, IDEAS'11 Proceedings 15th International Database Engineering & Applications Symposium, pages 247-248, Lisbon, 2011. Excellence in Research for Australia (ERA) indexed this symposium in the B category in its most recent hierarchy from 2010 [Era10]. (URL of the article: http://dl.acm.org/citation.cfm?id=2076623. 2076658&coll=DL&dl=ACM&CFID=63572418&CFTOKEN=57655636)
- Yann Gripay, Frédérique Laforest, François Lesueur, Nicolas Lumineau, Jean-Marc Petit, Vasile-Marian Scuturici, Samir Sebahi and Sabina Surdu, Colis-Track: Testbed for a Pervasive Environment Management System, EDBT 2012
   The 15th International Conference on Extending Database Technology, Berlin, 2012. The conference is indexed in the A category by ERA in 2010 [Era10]. (URL accepted papers: http://edbticdt2012.dima.tu-berlin.de/program/EDBT-papers/)
- Sabina Surdu, A new framework for evaluating performance in data stream monitoring applications with StreamInsight: StreamEval, MaCS 2012 Booklet of abstracts from The 9th Joint Conference on Mathematics and Computer Science (international conference), page 92, Siófok, 2012. (URL Booklet of abstracts: http://macs.elte.hu/downloads/abstracts/booklet.pdf)
- Sabina Surdu, A New Architecture Supporting The Sizing Window Effect With StreamInsight, Studia Universitatis Babeș-Bolyai Series Informatica,

LVI(4):111-120, 2011. The journal is indexed in the B+ category (BDI indexed) by CNCSIS in 2011 [CNC11].

- Sabina Surdu, Data stream management systems: a response to large scale scientific data requirements, Annals of the University of Craiova, Mathematics and Computer Science Series, 38(3):66-75, 2011. The journal is indexed in the B+ category (BDI indexed) by CNCSIS in 2011 [CNC11].
- Sabina Surdu, A new architecture for load shedding on data streams with StreamInsight: StreamShedder, University of Pitești Scientific Bulletin, Series Electronics and Computers Science, 11(2):57-64, 2011. The journal is indexed in the B+ category (BDI indexed) by CNCSIS in 2011 [CNC11].
- Sabina Surdu, A technique for constructing training sets in data stream mining: kSiEved Window Training Set, MDIS 2011 - Proceedings of The Second International Conference on Modelling and Development of Intelligent Systems, pages 180-191, Sibiu, 2011. (URL conference proceedings: http://conferences.ulbsibiu.ro/mdis/2011/Doc/Proceeding\_ mdis2011.pdf)
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- Sabina Surdu, Online political communication, Interdisciplinary New Media Studies Conference Proceedings (international conference), pages 55-58, Cluj-Napoca, 2009. (URL conference program: http://journalism. polito.ubbcluj.ro/inms/wp-content/uploads/2010/07/INMS\_ conference\_prog.pdf)

The following papers are under evaluation or about to be submitted to conferences or journals:

- **Sabina Surdu**, Yann Gripay, Jean-Marc Petit and Vasile-Marian Scuturici, paper sent to an A\* international conference 2012, under evaluation.
- Sabina Surdu, A new framework for evaluating performance in data stream monitoring applications with StreamInsight: StreamEval, Annales Universitatis Scientiarum Budapestinensis de Rolando Eötvös Nominatae Sectio Computatorica, 2012, under evaluation. The extended paper was sent together with the abstract having the same title, accepted at a previously mentioned international conference.
- Sabina Surdu și Vasile-Marian Scuturici, Assessing performance in data stream processing, to be submitted to IDEAS 2012 The 16th International Database Engineering & Applications Symposium, Prague, 2012. The symposium is indexed in the B category by ERA in 2010 [Era10].
- Sabina Surdu, Data stream processing: traditional vs. dedicated systems (SQL Server vs. StreamInsight), to be submitted to Studia Universitatis Babeş-Bolyai Series Informatica. The journal is indexed in the B+ category (BDI indexed) by CNCSIS in 2011 [CNC11].

#### **1** Thesis structure

In Chapter 1 we briefly describe the general problem of data stream processing using continuous queries and we present the ubiquitous network society, characterized by heterogeneous data processing in pervasive environments and applications. We present the challenge of optimizing queries when processing data streams from pervasive environments and we highlight the two main research directions we follow in this thesis: optimizing resource usage when processing queries on data streams and heterogeneous data management in pervasive application development. These applications contain static data, streams and functionalities [GLP10]. We discuss the original contributions from this thesis and we mention a list of papers, published in journals or presented at international conferences and published in proceedings. We also mention the papers we sent to conferences or journals, which are under evaluation or accepted for publication.

In Chapter 2 we present the state of the art in data stream processing. We introduce key data stream processing systems and discuss alternative approaches for query optimization, mostly oriented towards reducing system resource usage. We provide a comparative view on the discussed approaches.

In Chapter 3 we describe the techniques we propose in our thesis for optimizing resource usage in data stream processing. We analyse the sizing window effect, in order to determine an optimal window size for a query, so that resource usage is minimal and accuracy constraints are met. We discuss the kSiEved Window Training Set technique, a strategy for building training sets in data mining on data streams. This strategy aims at saving system resources and still meeting accuracy requirements.

In Chapter 4 we discuss the resource-aware architectures we designed, oriented towards reducing resource usage when processing continuous queries. StreamShedder and WindowSized are two such architectures for a Data Stream Management System, based on a commercial system for data stream processing. We briefly describe StreamEval, an application that evaluates performance variations when some conditions in the environment change. We also discuss SCIPE and InstantSchool-Know, two proposals we made for Data Stream Management Systems for specific application domains.

In Chapter 5 we move on to pervasive application development. We present the testbed we developed as part of a team, at LIRIS, INSA Lyon, for a system that manages pervasive environments, based on a scenario designed for such environments, in a medical context. This testbed can be used to assess pervasive application development. We present the design of a data-oriented pervasive application, using the SoCQ system [GFLP09], thereby homogeneously handling the data in the pervasive environment. We describe the application we implemented, which enables a user to write continuous queries, combining heterogeneous data.

In Chapter 6 we evaluate the development of pervasive applications, using multiple systems. We describe the proposed benchmark and conduct an experimental study. The results of the research presented in this chapter are part of a paper we submitted to an international conference and that is currently under evaluation.

In Chapter 7 we summarize the results we obtained on the chosen research directions: optimizing resource usage when processing queries on data streams and managing heterogeneous data in pervasive application development. We briefly discuss the techniques we developed for optimizing resource usage in data stream processing and the resource-aware architectures we designed to save resource usage when processing continuous queries on data streams. We touch on the testbed for pervasive application development and the benchmark for evaluating these applications. We describe future research directions driven by the results we obtained.

### 2 Processing data streams in pervasive environments

In the last couple of years we witnessed a shift in the traditional data processing paradigm, from the classical model, characterized by static data, to a dynamic model, which encompasses data characterized by high dynamics. In a growing number of fields, information takes the form of continuous streams. These are potentially unending sequences of data, which cannot be efficiently handled by traditional DBMSs [ACC<sup>+</sup>03]. Research teams from the academic world have developed prototypes to manage and process data streams. These are called Data Stream Management Systems (DSMSs). Industrial players also designed and launched fully featured DSMSs (a recent example is StreamInsight, released by Microsoft [KDA<sup>+</sup>10]).

In traditional databases, data has a static nature. It takes the form of finite, stored data sets, which are queried when necessary [ABB<sup>+</sup>04]. On the other hand, data streams are dynamic by definition. They are not permanently stored in a system. A query in this context executes in a perpetual manner on temporary data, which arrives at the system, is processed and in the end is discarded. A DSMS can execute a great number of complex continuous queries [ABB<sup>+</sup>03], which take into account multiple data streams. The stream data rate can greatly vary in time. Limited system resources must cope with these issues, when data processing must consider the temporal nature of the data.

In the applications from the the *ubiquitous network society* [Mur09], the user not only interacts with other users, but it also relates to objects from the environment, equipped with computational-enabled devices [Uni05]. In this context, data streams coexist with data modelled in different manners. Systems that manage such environments must consider static data, functionalities and data streams; pervasive environments are composed of such ingredients and are used to model the surrounding reality [GLL<sup>+</sup>12]. The integration of static data, data streams and functionalities querying capabilities in a unified, declarative model, opens other directions for query optimization in this new context, but similar to those found in traditional databases, based on SQL-like languages [Gri09]. Data streams and pervasive applications developed in the context of pervasive environments are the newcomers in the ubiquitous network society scenarios. Efficient resource management when processing data streams and the smooth development of pervasive applications are the necessary prerequisites in order to achieve the ubiquitous network society.

In this abstract we briefly describe the most important original contributions we present in our thesis.

## 3 Problem statement

In this thesis we investigate the general problem of query optimization, in the context of data stream processing in pervasive environments. We identify to main research directions, materialized in our published papers, listed in the preamble of this abstract:

- optimizing resource usage when processing queries on data streams;
- investigating heterogeneous data management in pervasive application development.

# 4 Optimizing resource usage when processing queries on data streams

One of the main problems that stream processing system designers are dealing with today is intensive resource usage. A system with limited resources needs to be able to handle a huge number of data sources, considerable data volumes, fast data rates and unpredictable spikes in data (like we show in [SS11]). The number of data sources, the data rates and the data volumes are constantly increasing. The data distribution can be variable and the system needs to be able to execute multiple complex queries, in a continuous manner [ABB<sup>+</sup>03]. In this context, we raise the following questions: how can the system function properly under these circumstances and how can the system performance be assessed?

In our thesis we describe the novel solutions we proposed to tackle this problem, presented on two research directions: (1) resource usage optimizing techniques in data stream processing and (2) designing resource-aware architectures for data stream processing, oriented towards saving system resources.

We enumerate the original contributions from our thesis, which aim at reducing resource usage when processing data streams.

#### 4.1 The sizing window effect

We develop the sizing window effect, an approach that aims at optimizing the resource usage (memory and CPU time), by computing an optimal window size for a given continuous query. We intend to improve this technique, so that the computation of the optimal window size can be performed by the system, in an automated manner. To the best of our knowledge, this is the first study that takes into account the size of the input window in order to reduce resource usage. We don't take into account windows that are semantically significant (for instance, a query that computes the average speed of cars on a road segment, in the last five minutes, needs a sliding window of fixed size). The semantics of these windows is derived from their temporal dimension. In our case, the semantics of the window is not connected to this parameter.

We briefly present the sizing window effect (in our thesis, we rigorously formalize the temporal domain, the data streams, the notion of query equivalence, the ideal result, the approximated result, the distance function and other concepts we use; in this abstract we concisely present them). A sliding window in this context is a contiguous portion of data from a data stream *S* [BBD<sup>+</sup>02]. If its temporal boundaries are instants  $t_i$  and  $t_j$ , we will denote this window by  $SW_{ij}(S)$ .

Let Q be a query whose execution produces a stream of aggregate results over time.  $t_c \in T$  is the current timestamp, where T is the chosen temporal domain.  $t_i \in T$  is a timestamp that marks the starting point of a window in time and  $t_0$  is the timestamp of the first emitted element on stream S. Initially  $t_i = t_c$ . We denote by CrtTS the set of all timestamps from T, which will be assigned to  $t_c$ . We go through the following stages:

- 1. We establish an accuracy threshold  $\epsilon$ . In order to obtain equivalent queries and valid answers, the difference between the ideal results and the approximated results must not exceed the accuracy threshold.
- We compute the ideal result R<sub>sc</sub> of query Q executed on data stream S, at current time instant t<sub>c</sub>:

$$R_{s_c} = Q(S, t_c) = Q(SW_{0c}(S), t_c), R_{s_c} \in \mathbb{R}$$
(1.1)

where  $SW_{0c}(S)$  is a sliding window from the data stream *S*, whose temporal boundaries are  $t_0$  and  $t_c$ . We say this result is *ideal*, because it takes into account all the elements arrived on the stream until the current time instant.

We constantly decrease t<sub>i</sub>. We compute the approximated results R<sub>w<sub>cσic</sub></sub> of query Q executed on the sliding windows SW<sub>ic</sub>(S), at current timestamp t<sub>c</sub>. For each temporal value t<sub>i</sub>, the window dimension SW<sub>ic</sub>(S) is σ<sub>ic</sub>, representing the number of temporal instants contained by the window:

$$R_{w_{c_{\sigma_{ic}}}} = Q(SW_{ic}(S), t_c), R_{w_{c_{\sigma_{ic}}}} \in \mathbb{R}$$

$$(1.2)$$

We compute the distances between the ideal result and the approximated result, using a distance function, for each value of the  $t_i$  timestamp:

$$distance_{agg_{c_{\sigma_{ic}}}}(R_{s_c}, R_{w_{c_{\sigma_{ic}}}}) = \mid R_{s_c} - R_{w_{c_{\sigma_{ic}}}} \mid$$
(1.3)

- 4. We repeat steps 2 and 3 for all the values the current timestamp  $t_c$  takes from *CrtTS*.
- 5. After finalizing step 4 (when  $t_c$  has taken all the values from CrtTS), we compute the average of the distances between the approximated results and the ideal results over time, for each window dimension  $\sigma_{ic}$ :

$$AvgDistance(\sigma_{ic}) = \frac{\sum_{t_c \in CrtTS} distance_{agg_{c_{\sigma_{ic}}}}(R_{s_c}, R_{w_{c_{\sigma_{ic}}}})}{|CrtTS|}$$
(1.4)

The optimal window dimension for query Q is the smallest window dimension for which the average distance between the approximated results and the ideal results is less than the accuracy threshold  $\epsilon$ .

In our experiments, conducted on a set of aggregate queries, we use the data from the Linear Road benchmark [ACG<sup>+</sup>04]. This is simulated data, related to road traffic on expressways, which are all divided into 100 segments.

For each of the following aggregate queries we will apply the previously described tactic.

Query 1: Calculate the average number of cars per time unit which have been travelling on a given segment.

Query 2: Calculate the average speed on a given segment.

Query 3: Calculate the average toll paid by a car (on all the segments).

In the case of Query 1 (figure 1.1, which we published in [Sur11a]), we notice that the average distance between the ideal results and the approximated results is less than the accuracy threshold 1, for all the window dimensions greater than 1000 time instants.



Figure 1.1: Average number of cars per time unit, separately computed for 5 segments.



Figure 1.2: Average speed, separately computed for five segments



Figure 1.3: Average toll paid by 5 cars (computed separately for every car)

In the case of Query 2 (figure 1.2), the average distance between the ideal results and the approximated results is less than the accuracy threshold 1, for all the window dimensions greater than 10000 time instants.

In the case of Query 3 (figure 1.3), the average distance between the ideal results and the approximated results is less than the accuracy threshold 0.1, for all the window dimensions greater than 6000 time instants.

The administrator of an application that implements Linear Road can establish an accuracy threshold for Query 1 on the output data (from the query executed on sliding windows), so that it doesn't differ by more than 1 from the ideal result. The system can execute this query on a window of 1000 time instants. Similar constraints can be formulated for the other queries as well. The results of this research are published in [SS11].

#### 4.2 kSiEved Window Training Set

One of the challenges encountered in data mining today is applying techniques specific to this process on continuous data streams [ZB03]. We develop a technique that takes into account system resources when constructing training sets for data mining algorithms on data streams: the kSiEved Window Training Set (kSEWT). This is the first method that "sieves" a data stream depending on some parameters, to construct training sets in this context, while meeting accuracy requirements. We define a new data model, the kSiEved model, which is based on kSiEved windows, built from sliding windows by applying functions that extract positions from a window. We rigorously define this model in our thesis.

kSEWT computes correct results, on sliding windows  $SW_{ic}$ , at every time instant  $t_c$  (we omit the stream *S* in the definition of these windows to simplify the notations). For each such window, kSEWT builds kSiEved windows  $SEW_{ic}(k)$ , based on a parameter *k*, which varies in time. This parameter generates a "sieve" with holes, which will "sieve" the elements from the window  $SW_{ic}$ , yielding the kSiEved window  $SEW_{ic}(k)$ . On this window we compute queries results as well. kSEWT estimates the accuracy of the obtained results on kSiEved windows when compared to the correct results by using a distance function. Depending on the average of the computed distances, we choose parameter *k* (its greatest value), for which the average of the distances from the correct result does not exceed a previously established error threshold  $\delta$ . Parameter *k* provides the kSiEved Window Training Set, which contains all the kSiEved windows with parameter k obtained in the experiment.

We present the experimental results we obtained on a data set with a uniform distribution. By applying kSEWT we obtained the graph from figure 1.4. If we choose a threshold  $\delta = 0.5$ , from this graph we notice we can apply "sieves" with parameter k = 2, when building the training set. This means we discard half of the input tuples, while still meeting formulated accuracy requirements. This way we are reducing resource usage in a significant manner. This research is published in

[Sur11b].



Figure 1.4: Average distance between the correct results and the results of queries executed on kSiEved windows

#### 4.3 WindowSized

Deepening our study on the sizing window effect, we propose a new resourceaware architecture to implement this effect, using Microsoft StreamInsight [KDA<sup>+</sup>10]: WindowSized. The main contribution of this architecture is the integration of the WindowSizing module in a monitoring application developed with StreamInsight. WindowSizing interacts with the query engine, with the data sources interfaces - to modify the size of the window and with the output devices interfaces - to obtain the queries results.

Figure 1.5 depicts the main components of such an architecture. The lower level from the architecture (containing *Event sources, Input adapters, StreamInsight query engine, Output adapters* and *Event targets*) is taken from the architecture proposed by Microsoft, in order to implement an application with StreamInsight [SIA]. WindowSized is based on a tipical StreamInsight application design priciples, with the following elements: data sources, input adapters, continuous queries on the server,

output adapters and data consumers [GSK<sup>+</sup>09]. Our contribution is represented by the integration of the WindowSizing module in such an architecture. The obtained results are published in [Sur11a].



Figure 1.5: The WindowSized architecture

#### 4.4 StreamShedder

We develop the StreamingTraffic traffic monitoring application. We propose a new architecture for such a monitoring application implemented with the StreamInsight platform [KDA<sup>+</sup>10]. We develop the load shedding [ABB<sup>+</sup>04] module Stream-Shedder and we recommend its integration in the architecture of StreamingTraffic. StreamShedder performs data elimination operations in a parameterized manner, taking into account system resources and query latency. The obtained architecture integrates load shedding strategies with a commercial stream processing system to obtain superior performances when processing continuous queries.

Figure 1.6 depicts the modified architecture of a monitoring application implemented with StreamInsight, which encompasses the StreamShedder module. Like in the case of WindowSized, the lower level from the architecture (with *Event sources, Input adapters, StreamInsight query engine, Output adapters* and *Event targets*) is taken from the architecture proposed by Microsoft, in order to implement an application with StreamInsight [SIA]. Our contribution is represented by the integration of the StreamShedder module in such an architecture. We will call this enriched architecture StreamShedder as well, based on the module that performs the load shedding operations.



Figure 1.6: The StreamShedder architecture

StreamShedder is a software module implemented in C#. It communicates with a memory and processor monitoring module (CPU and Memory), which provides the used system resources. Depending on the user specified memory and CPU time thresholds, StreamShedder can command the input adapters to eliminate some tuples. StreamShedder also monitors the latency of the queries on the server. Based on the data it receives, it "tells" the input adapters what tuples to eliminate. To assess the impact on the application semantics, this module also communicates with output adapters, obtaining queries results. The results of this research are published in [Sur11d].

#### 4.5 StreamEval

We develop a solution that evaluates performance variations when different conditions from the environment change (for instance, when we manipulate the data rate of the data sources that feed the continuous queries on the server): StreamEval. In the implementation of the monitoring application and of our proposed framework we use the previously mentioned commercial platform developed by Microsoft in the last couple of years: StreamInsight [AGR<sup>+</sup>09]. We use the StreamingTraffic monitoring application we developed for the StreamShedder architecture (section 4.4). We denote by *DR* the data rate of the input stream, defined as the number of elements that arrive on the input stream *S* every second. We will use this notation when we change the data source's data rate. We use the (slightly changed) notation *ConsumedGate* from [Mon] to refer to the point immediately following the input adapters (at the first operator from a continuous query *Q*).

We use the query monitoring attributes provided by the ManagementService API [Mon]. Like in [Mon], we are interested in monitoring the average consumed latency, between two time instants  $t_1$  and  $t_2$ . Therefore, we evaluate the number of processed tuples *TupleCount* and the response time *Latency* at *ConsumedGate*, at moments  $t_1$  and  $t_2$ . We denote the average consumed latency by *AvgLat*. We compute *AvgLat* by applying the formula from [Mon]:

$$AvgLat = (Latency_{t_2} - Latency_{t_1}) / (TupleCount_{t_2} - TupleCount_{t_1}).$$
(1.5)

We change the data rate of the data source as follows. We start with value 1 for *DR* (one event every second) and we evaluate the corresponding *AvgLat* value. We increase *DR* up to 500 events per second. The metric *AvgLat* stays under one millisecond. Even for values of 1000 events / second for *DR*, which exceed the requirements of StreamingTraffic, *AvgLat* maintains around the same value. This research is briefly described in the abstract [Sur12a]. The extended paper is under evaluation [Sur12b].

#### 4.6 SCIPE

We develop a set of principles, SCIPE (*SCIentific data stream processing PrinciplEs*), oriented towards the realization of a Sci-DSMS, a Data Stream Management System that handles very large data from domains connected with exact sciences. Research communities from exact sciences have been working with data in the order of petabyes and this data is expected to exceed exabytes dimensions in the years to come [BLW09]. In this context, we investigate the possibility of implementing a DSMS, which answers the needs of exact sciences communities. Considering the objectives

of a particular domain can lead to optimizing resource usage in continuous queries on data streams.

We describe SCIPE:

1. When possible, the system should process and subsequently summarize or eliminate arriving elements. This principle has a significant impact on memory usage, by keeping elements in the form of a summary, if they are required for future processing.

2. If each element needs to be individually stored, then only the elements from the recent past should be kept and old elements should be eliminated or summarized.

3. The designed system should encompass elements revision possibilities (this strategy is borrowed from [AAB<sup>+</sup>05]).

4. Load shedding should be performed in a semantic manner, dependent on the application domain (Aurora [ACC<sup>+</sup>03] is a system that performs semantic load shedding).

5. Queries should be written in a user-friendly manner, combining visual languages and a declarative SQL interface (we combine the approaches from [ACC<sup>+</sup>03] and [ABW06]).

SCIPE and its motivation are published in [Sur11c].

#### 4.7 InstantSchoolKnow

We analyse the educational field and the ways in which data streams can optimize the educational processes. We develop EdStream, a set of rules that can be applied when implementing an educational monitoring platform based on data stream processing. We propose the design of an educational monitoring platform, Instant-SchoolKnow. Its purpose is to continuously acquire data from educational institutions (registered on the platform), to analyse this data using a continuous processing paradigm and to publish the results of this analysis in real time. To achieve this goal the following stages must be accomplished: the registration on the Instant-SchoolKnow platform, data acquisition, data analysis and data publishing. Unlike current approaches, InstantSchoolKnow aims at unifying e-learning and students monitoring capabilities in a single platform. This research is published in [Sur11e].

### 4.8 A platform for accessing data on smart mobile decives

We propose an architecture for the implementation of an online platform with content oriented towards political communication. In its initial stage the data has a static nature and can be accessed from smart mobile devices. We intend to extend this new media platform with stream and services processing capabilities, in the context of a pervasive environment. The research is published in [Sur09].

# 5 Managing heterogeneous data in pervasive application development

A considerable number of scenarios and pervasive applications based on these scenarios contain static data (similar to data from relational databases), data streams and distributed functionalities or services [GLP10], according to the real, everyday life situations they model. In order to manage all these elements from a pervasive environment, ad hoc programming is usually used, which integrates multiple programming models: imperative and declarative languages and network protocols [Gri09]. Solutions developed in this manner are complex in their implementation and the development time is high. We investigate alternative approaches to pervasive application development and methods to assess the development process.

We enumerate the original contributions from our thesis, which we obtained in the context of heterogeneous data management in pervasive applications.

#### 5.1 Heterogeneous data management in a pervasive environment

We tackle one of the main challenges in pervasive computing: facilitating pervasive application development. We describe a scenario for container monitoring in a medical context, which refers to the transportation of medical content in containers equipped with sensors. Based on this scenario, we discuss a testbed, which can be used when developing applications and assessing the development process and we show how to build a pervasive application using the SoCQ (Service-oriented Continuous Query) system [GFLP09]. The scenario, its simulation as a testbed, its visualization and the developed pervasive application represent the intrinsic contributions of this research, which we developed together with the team we worked with, at LIRIS, INSA Lyon. The results of this research are presented in a paper accepted at an international conference and will be published [GLL<sup>+</sup>12].

To interact with the query engine, we implement a Web ASP.NET application. This application allows a developer to write continuous queries, which mix hete-

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Longitude REAL				
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SPONSE:	LocationTimestamp	Latitude	Longitude	SCHEMA
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Figure 1.7: The Web application that allows continuous query writing

rogeneous data from the environment, by using a query language similar to SQL, specific to the SoCQ system [Gri09]. If we want to monitor the location of each car at each instant, we write a query in this language, which generates all the car locations as a result, in a real time manner. Figure 1.7 depicts the Web application, a query and its results.

#### 5.2 AgilBench

We propose a benchmark for evaluating pervasive application development. We use multiple systems for this purpose and we conduct an experimental study. The results of this research were included in a paper we submitted to an international conference and that is currently under evaluation [SGPS12].

#### 6 Conclusions and future research directions

The two main research directions we followed led to developing architectures, strategies and techniques to optimize resource usage when processing data streams, but also to the implementation of a testbed and a benchmark in the context of pervasive application development. These contributions were published in journals or proceedings of international conferences. One additional paper is accepted for publication and two other papers are currently under evaluation.

The fields of data streams and pervasive applications are advancing on a daily basis. We can certainly expect that our proposals will suffer changes in time. We intend to automate the sizing window effect, so that the system can automatically choose the optimal window size and to enhance the resource-aware architectures we proposed, so that all the decisions are taken by the system, with no user intervention. We want to add new services in the testbed and enrich the benchmark we implemented to evaluate pervasive application development. We assess the possibility of designing a system, capable of managing pervasive environments, which allows the total replacement of the scenario that models a pervasive environment, which no changes in its implementation. The most powerful systems (like SoCQ) need new modules if the data access mechanisms change, once the scenario is replaced.

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